# Dendritic Computing in the Lattice Domain 

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## Overview

Part I: Background

- Rationale For Lattice Based Dendritic Computing
- Lattice Neural Networks (LNNs)
- Problems Associated With Current LNNs

Part II: A Novel Two Metric Model

- The Basic Idea Behind The Two Metric Model
- Lattice Metrics and Hyperplanes
- The Two Metric Model and Examples
- Concluding Remarks and Questions


## Rationale for Dendritic Computing

- Basic Goal: A return of ANNs to its Roots in Neurobiology and Neurophysics
- Radial Basis Function NNs, SVM, Boltzmann Machines, etc., bear little resemblance to biological neural networks
- Dendrites make up more than $50 \%$ of a neuron's membrane
- Dendrites make up the largest component in both surface area and volume of the brain
- A neuron in the cortex typically sends messages to approximately $10^{4}$ other neurons.


## Rationale for Dendritic Computing

- Dendrites and dendritic spines are major postsynaptic targets of presynaptic inputs
- The number of synapses on a single neuron ranges between 500 and 200,000
- The number of synapses in the human brain ranges between 60 trillion and 240 trillion $\left(240 \times 10^{12}\right)$
- These synapses reside on 10 to 20 billion neurons


## Rationale for Dendritic Computing

- Recent research results demonstrate that the dynamic interaction of inputs in dendrites containing voltage-sensitive ion channels make them capable of realizing nonlinear interactions, logical operations, and possibly other local domain computation (Poggio, Koch, Shepherd, Rall, Segev, Perkel, et.al.)
- Based on their experimentations, these researchers make the case that it is the dendrites and not the neural cell bodies that are the basic computational units of the brain.
- Thus, when attempting to model artificial brain networks, one cannot ignore dendrites


## Rationale for Lattice Computing

- Neurons with dendrites can function as many independent subunits with each unit being able to implement a rich repertoire of logical operations
- Logical functions such as XOR, AND, OR, and NOT; Koch, Riesenhuber, Poggio, Setiono, Segev and others
- Lattce operations involve only max, min, and addition; i.e., $\vee, \wedge$, and +
- Thus, lattice operations provide for extremely fast neural computation and fast learning methods

Biological Neurons


Figure 1: Simplified sketch of the processes of a biological neuron.

## Dendritic Computation: Assumptions

- The postsynaptic neuron $M_{j}$ receives input from $n$ presynaptic neurons $N_{1}, \ldots, N_{n}$.
- Each input neuron $N_{i}$ has axonal branches that terminate at various synaptic regions of $M_{j}$.
- The synaptic regions are distributed along a finite number of dendrites $d_{1}, \ldots, d_{K(j)}$.
- Incoming information from axonal branches is transformed in the synaptic interaction
- The transformed data will result in either an excitatory postsynaptic response or an inhibitory postsynaptic response in the dendrites membrane.


## Postsynaptic neuron with dendritic structures



Terminal branches of axonal fibers originating from the presynaptic neurons make contact with synaptic sites on dendritic branches of $M_{j}$

## Dendritic Computation: Mathematical Model

The computation performed by the $k$ th dendrite for input $\mathbf{x}=\left(x_{1}, \ldots, x_{n}\right)^{\prime} \in \mathbb{R}^{n}$ is given by

$$
\tau_{k}^{j}(\mathbf{x})=p_{j k} \bigwedge_{i \in I(k)} \bigwedge_{\ell \in L(i)}(-1)^{1-\ell}\left(x_{i}+w_{i j k}^{\ell}\right)
$$

where

- $x_{i}$ - value of neuron $N_{i}$;
- $I(k) \subseteq\{1, \ldots, n\}-$ set of all input neurons with terminal fibers that synapse on dendrite $d_{j k}$;
- $L(i) \subseteq\{0,1\}$ - set of terminal fibers of $N_{i}$ that synapse on dendrite $d_{j k}$;
- $p_{j k} \in\{-1,1\}-$ EPSP/IPSP of $d_{j k}$.


## Dendritic Computation: Mathematical Model

- The value $\tau_{k}^{j}(\mathbf{x})$ is passed to the cell body and the state of $M_{j}$ is a function of the input received from all its dendritic postsynaptic results. The total value received by $M_{j}$ is given by

$$
\tau^{j}(\mathbf{x})=p_{j} \bigwedge_{k=1}^{K(j)} \tau_{k}^{j}(\mathbf{x})
$$

## The Capabilities of an SLLP

- An SLLP can distiguish between any given number of pattern classes to within any desired degree of $\varepsilon>0$.
- More precisely, suppose $X_{1}, X_{2}, \ldots, X_{m}$ denotes a collection of disjoint compact subsets of $\mathbb{R}^{n}$.
- For each $p \in\{1, \ldots, m\}$, define $Y_{p}=\bigcup_{j=1, j \neq p}^{m} X_{j}$
$\varepsilon_{p}=\mathrm{d}\left(X_{p}, Y_{p}\right)>0$
$\varepsilon_{0}=\frac{1}{2} \min \left\{\varepsilon_{1}, \ldots, \varepsilon_{m}\right\}$.
- As the following theorem shows, a given pattern $\mathrm{x} \in \mathbb{R}^{n}$ will be recognised correctly as belonging to class $C_{p}$ whenever $\mathbf{x} \in X_{p}$


## The Capabilities of an SLLP

- Theorem. If $\left\{X_{1}, X_{2}, \ldots, X_{m}\right\}$ is a collection of disjoint compact subsets of $\mathbb{R}^{n}$ and $\varepsilon$ a positive number with $\varepsilon<\varepsilon_{0}$, then there exists a single layer lattice based perceptron that assigns each point $\mathrm{x} \in \mathbb{R}^{n}$ to class $C_{j}$ whenever $\mathrm{x} \in X_{j}$ and $j \in\{1, \ldots, m\}$, and to class $C_{0}=\neg \bigcup_{j=1}^{m} C_{j}$ whenever $\mathrm{d}\left(\mathrm{x}, X_{i}\right)>\varepsilon, \forall i=1, \ldots, m$. Furthermore, no point $\mathrm{x} \in \mathbb{R}^{n}$ is assigned to more than one class.


## Graphical Interpretation of Theorem



Any point in the set $X_{j}$ is identified with class $C_{j}$; points within the $\epsilon$-band may or may not be classified as belonging to $C_{j}$, points outside the $\epsilon$-bands will not be associated with a class $C_{j} \forall j$.

## Learning in LNNs

- Early training methods were based on the proofs of the preceding Theorems.
- All training algorithms involve the growth of axonal branches, computation of branch weights, creation of dendrites, and synapses.
- The first training algorithm developed was based on elimination of foreign patterns from a given training set (min or intersection).
- The second training algorithm was based on small region merging (max or union).



Left : Two class data set. Right : The elimination method.



## Left : The merging method. Right : Boundary readjustment.

## SLLP Using Elemination VS MLP


(a) SLLP: 3 dendrites, 9 axonal branches. (b) MLP 13 hidden neorons and 2000 epochs.

## SLLP Using Merging



During training, the SLLP grows 20 dendrites, 19 excitatory and 1 inhibitory (dashed).
(2)

## Another Merging Example



Learning in LNNs

| Classifier | Recognition |
| :--- | :---: |
| SLLP (elimination) | $98.0 \%$ |
| Backpropagation | $96 \%$ |
| Resilient Backpropagation | $96.2 \%$ |
| Bayesian Classifier | $96.8 \%$ |
| Fuzzy LNN | $100 \%$ |

UC Irvine Ionosphere data set (2-class problem in $\mathbb{R}^{34}$ with training set $=65 \%$ of data set)

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## Learning in LNNs

| Classifier | Recognition |
| :--- | :---: |
| Fuzzy SLLP (merge/elimination) | $98.7 \%$ |
| Backpropagation | $95.2 \%$ |
| Fuzzy Min-Max NN | $97.3 \%$ |
| Bayesian Classifier | $97.3 \%$ |
| Fisher Ratios Discimination | $96.0 \%$ |
| Ho-Kashyap | $97.3 \%$ |

Fisher's Iris Data Set. A 3-class problem in $\mathbb{R}^{4}$ with training set $=50 \%$ of data set.

## Dendritic Model of an Associative Memory



Topology of the dendritic associative memory based on the dendritic model. The network is fully connected.

## Patterns to Store



Top row represents the patterns $\mathrm{x}^{1}, \mathrm{x}^{2}$, and $\mathrm{x}^{3}$, while the bottom row depicts the associated patterns $\mathbf{y}^{1}, \mathbf{y}^{2}$, and $\mathbf{y}^{3}$. Here $n=2500$ and $m=1500$.

## Successful Recall of Noisy Patterns



The top row shows noisy input patterns. Bottom row shows perfect racall association.

## Problems with the Hyperbox Approach




The triangular data can never be modeled exactly using either elimination or merging.

## Learning in LNNs

- A. Barmpoutis extended the elimination method to arbitrary orthonormal basis settings.
- Basic Equation:

$$
\tau_{k}^{j}(\mathbf{x})=p_{k}^{j} \bigwedge_{i \in I(k)} \bigwedge_{\ell \in L(i)}(-1)^{1-\ell}\left(\left(\mathbf{R}_{k} \cdot \mathbf{x}\right)_{i}+w_{i j k}^{\ell}\right)
$$

## Problems with Rotations



The minimal standard $L_{\infty}$-rectangle and the minimal $45^{\circ} \mathrm{OB}$-rectangle are as shown.

## Lattice Metrics

- $L_{1}$ metric $d_{1}(\mathbf{x}, \mathbf{y})=\sum_{i=1}^{n}\left|x_{i}-y_{i}\right|$
- $L_{\infty}$ metric $d_{\infty}(\mathbf{x}, \mathbf{y})=\bigvee_{i=1}^{n}\left|x_{i}-y_{i}\right|$
- Hausdorff metric based on either the $d_{1}$ or $d_{\infty}$ metric


The $d_{1}$ Sphere in $\mathbb{R}^{3}$

## Hyperplanes

- A hyperplane in $\mathbb{R}^{n}$ is defined by

$$
a_{1} x_{1}+a_{2} x_{2}+\cdots+a_{n} x_{n}=b,
$$

where not all the $a_{i}$ 's are zero

- If $a_{i} \in\{-1,1\} \forall i$, then the hyperplane is an $L_{1}$-hyperplane
- If $a_{i}= \pm 1$ and $a_{j}=0 \forall j \neq i$, then the hyperplane is an $L_{\infty}$-hyperplane.


## Pertinent Hyperplane Properties

- A hyperplane can also be defined by the function

$$
f(\mathbf{x})=a_{1} x_{1}+a_{2} x_{2}+\cdots+a_{n} x_{n}-b=0
$$

- A hyperplane separates $\mathbb{R}^{n}$ into two half-spaces $H^{+}$(i.e. $f(\mathbf{x}) \geq 0$ ) and $H^{-}$(i.e. $f(\mathbf{x}) \leq 0$ )
- Up to parallelism, there are $n L_{\infty}$-hyperplanes and $2^{n-1} L_{1}$-hyperplanes in $\mathbb{R}^{n}$


## Summation Formulae for $L_{1}$-Hyperplanes

Starting with the summations

$$
\begin{aligned}
& L_{1}^{1}(\mathbf{x})=x_{1}+x_{2} \\
& L_{2}^{1}(\mathbf{x})=-x_{1}+x_{2}
\end{aligned}
$$

it is easy to generate all summations for a given dimension $n$ using the recursion formula:

$$
L_{j}^{n-1}(\mathbf{x})=\left\{\begin{array}{lll}
L_{j}^{n-2}(\mathbf{x})+x_{n} & \text { if } & j=1, \ldots, 2^{n-2} \\
-L_{i}^{n-2}(\mathbf{x})+x_{n} & \text { if } & j=2^{n-2}+i,
\end{array}\right.
$$

where $i=1, \ldots, 2^{n-2}$.


Point $\mathbf{x}=\left(x_{1}, x_{2}\right)$ is in the triangle $\Leftrightarrow$

$$
\begin{aligned}
& L_{1}(\mathbf{x}) \wedge\left(4-L_{1}(\mathbf{x})\right) \wedge\left(L_{2}(\mathbf{x})+2\right) \wedge-L_{2}(\mathbf{x}) \\
& \quad \wedge x_{1} \wedge\left(2-x_{1}\right) \wedge x_{2} \wedge\left(2-x_{2}\right) \geq 0
\end{aligned}
$$

## The Two Metric Model

- The previous inequality is equivalent to

$$
\begin{aligned}
\tau(\mathbf{x}) & =\left[\bigwedge_{i=1}^{2} \bigwedge_{\ell=0}^{1}(-1)^{1-\ell}\left(L_{i}(\mathbf{x})+\omega_{i}^{\ell}\right)\right] \\
& \wedge\left[\bigwedge_{i=1}^{2} \bigwedge_{\ell=0}^{1}(-1)^{1-\ell}\left(x_{i}+w_{i}^{\ell}\right)\right] \geq 0
\end{aligned}
$$

where $\omega_{1}^{1}=\omega_{2}^{0}=w_{1}^{1}=w_{2}^{1}=0, \omega_{2}^{0}=-4$, $\omega_{2}^{1}=2$, and $w_{1}^{2}=-2=w_{2}^{0}$.

## The Two Metric Model

- This generalizes to

$$
\begin{aligned}
\tau_{k}^{j}(\mathbf{x}) & =\left[p_{k}^{j} \bigwedge_{i \in I(k)} \bigwedge_{\ell \in L(i)}(-1)^{1-\ell}\left(x_{i}+w_{i j k}^{\ell}\right)\right] \\
& \wedge\left[q_{k}^{j} \bigwedge_{i \in J(k)} \bigwedge_{\ell \in L^{\prime}(i)}(-1)^{1-\ell}\left(L_{i}(\mathbf{x})+\omega_{i j k}^{\ell}\right)\right]
\end{aligned}
$$

where $\omega_{h j k}^{\ell}$ is the synaptic weight at synapse of $L_{i}, J(k) \subseteq\left\{1, \ldots, 2^{n-1}\right\}$ set of all input neurons $L_{i}$ with terminal fibers on $d_{j k}$.


Two metric error: $\operatorname{Area}\left(P^{2} \cap H^{2}\right)-\pi r^{2}$, VS single metric error: $\operatorname{Area}\left(H^{2}\right)-\pi r^{2}=r^{2}(4-\pi)$.


Left : LNN solving the triangle problem derived from learning (elimination). Right : LNN after Pruning.


## Left : The 2-D XOR Problem. Right : LNN derived from learning followed by pruning



Left : The 2-D XOR Problem. Right : LNN derived from learning followed by pruning
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## Questions?

## Thank you!

